

## Exercise 2: Flirting with Synchrony and Asynchrony

### Task 1: Growing Balls

Denote by  $B(v, r)$  the ball of radius  $r$  around  $v$ , i.e.,  $B(v, r) = \{u \in V : \text{dist}(u, v) \leq r\}$ . Consider the following partitioning algorithm.

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**Algorithm 1** Cluster construction.  $\rho \geq 2$  is a given parameter.

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1: while there are unprocessed nodes do
2:   select an arbitrary unprocessed node  $v$ ;
3:    $r := 0$ ;
4:   while  $|B(v, r + 1)| > \rho|B(v, r)|$  do
5:      $r := r + 1$ 
6:   end while
7:   makeCluster( $B(v, r)$ )           // all nodes in  $B(v, r)$  are now processed
8:   remove all cluster nodes from the current graph
9: end while
10: select intercluster edges
```

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- Show that Algorithm 1 constructs clusters of radius at most  $\log_\rho n$ .
- Show that Algorithm 1 produces at most  $\rho n$  intercluster edges.
- For  $k \in \{1, \dots, \lceil \log n \rceil\}$ , determine an appropriate choice  $\rho(k)$ , proving the precondition of Corollary 2.14!

### Task 2: Showing Dijkstra, and Bellman & Ford the Ropes

- Show that if the asynchronous Bellman-Ford algorithm from the lecture is executed synchronously, it sends only  $\mathcal{O}(|E|)$  messages.
- Use this to construct an asynchronous BFS tree construction algorithm of time complexity  $\mathcal{O}(D)$  that uses  $\mathcal{O}(|E|D)$  messages and terminates. You may assume that  $D$  is known here.
- Can you give an asynchronous Bellman-Ford-based algorithm that sends  $\mathcal{O}(|E| + nD)$  messages and runs for  $\mathcal{O}(D^2)$  rounds?

**Hint:** Either answer is feasible, provided it is backed up by appropriate reasoning!

### Task 3\*: Liaison with Leslie Lamport

- Look up what Lamport causality, Lamport clocks, and Lamport vector clocks are.
- Contemplate their relation to synchronizers and what you've learned in the lecture.
- Discuss your findings in the exercise session!