Exercise 8: Don't get Lost

Task 1: . . . everything is (probably) going to be fine

An event occurs *with high probability (w.h.p.),* if its probability is, for any choice of $c \in \mathbb{R}_{\geq 1}$, at least $1 - n^{-c}$. Here n is the input size (in our case, $n = |V|$), and c is a (user-provided) parameter, very much like the ϵ in a $(1 + \epsilon)$-approximation algorithm.

Algorithm 1 Code for generating a random ID at node v.

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1: id_v ← ⌈c log n⌉ random bits
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a) Suppose that some algorithm $\mathcal{A}$ is called ten times, and each call succeeds w.h.p. Pick c such that for $n \geq 10$, all ten calls of $\mathcal{A}$ all succeed with a probability of at least 0.999.

   Hint: Union bound.

b) Let $\mathcal{E}_1, \ldots, \mathcal{E}_k$ be polynomially many events, i.e., $k \in n^{\mathcal{O}(1)}$, each of them occurring w.h.p. Show that $\mathcal{E} := \mathcal{E}_1 \cap \cdots \cap \mathcal{E}_k$, the event that all $\mathcal{E}_i$ happen, occurs w.h.p.

c) Consider Algorithm 1, which generates random node IDs. Fix two distinct nodes $v, w \in V$ and show that w.h.p., they have different IDs.

d) Show that w.h.p., Algorithm 1 generates pairwise distinct node IDs.

Task 2: . . . in the Steiner Forest!

In this exercise, we're going to find a 2-approximation for the Steiner Tree problem on a weighted graph $G = (V, E, W)$, as defined in an earlier exercise; we use the Congest model. Denote by T the set of nodes that need to be connected, and by $G_T = (T, \binom{T}{2}, W_T)$ the terminal graph.

a) For each node v, denote by t_v the closest node in T. Show that all $v \in V$ can determine t_v along with the weighted distance $\text{dist}(v, t_v)$ in

$$\max_{v \in V}\{\text{hop}(v, t_v)\} + \mathcal{O}(D)$$

   rounds,[1] where $\text{hop}(v, t_v)$ denotes the length of the minimum hop path from v to t_v.

   Hint: This essentially is a single-source Moore-Bellman-Ford with a virtual source connected to all nodes in T.

b) Consider a terminal graph edge $\{t_v, t_w\}$ "witnessed" by G-neighbors v and w with $t_v \neq t_w$, i.e., v and w know that $\text{dist}(t_v, t_w) \leq \text{dist}(t_v, v) + W(v, w) + \text{dist}(w, t_w)$. Show that if there are no such v and w with $\text{dist}(t_v, t_w) = \text{dist}(v, t_v) + W(v, w) + \text{dist}(w, t_w)$, then $\{t_v, t_w\}$ is not in the MST of G_T!

   Hint: Observe that G is partitioned into Voronoi cells $V_t = \{v \in V \mid t_v = t\}$, and that in the above case any shortest t_v-t_w path must contain a node u with $t_u \notin \{t_v, t_w\}$, i.e., cross a third Voronoi cell. Conclude that $\{t_v, t_w\}$ is the heaviest edge in the cycle (t_v, t_u, t_w, t_v).

[1] These are partial shortest-path trees rooted in each $t \in T$.

c) Show that an MST of G_T can be determined and made globally known in $\mathcal{O}(|T|+D)$ additional rounds.

Hint: Use the distributed variant of Kruskal's algorithm from the lecture.

d) Show how to construct a Steiner Tree of G of at most the same weight as the MST of the terminal graph in additional $\max_{v \in V}\{\text{hop}(v, t_v)\}$ rounds.

Hint: Modify the previous step so that the "detecting" pair v, w with $\text{dist}(t_v, t_w) = \text{dist}(v, t_v) + W(v, w) + \text{dist}(w, t_w)$ is remembered. Then mark the respective edges $\{v, w\}$ and the leaf-root-paths from v to t_v and w to t_w for inclusion in the Steiner Tree.

e) Conclude that the result is a 2-approximate Steiner Tree. What is the running time of the algorithm?

Hint: Recall Task 2 from Exercise 6.

Task 3*: ...under a Heap of Presents

weight	RGB
1	$(255, 255, 0)$
2	$(34, 139, 34)$
3	$(165, 42, 42)$
5	$(255, 0, 0)$
20	$(193, 255, 244)$

a) Determine an MST of the graph given in Figure 1!

b) Color each MST edge. The edge colors are given in the table above, i.e., an edge of weight 1 has color $(255, 255, 0)$.

c) Look for other Christmas trees in the computer science literature!

Hint: xkcd.

d) Have a Merry Christmas and a Happy New Year!

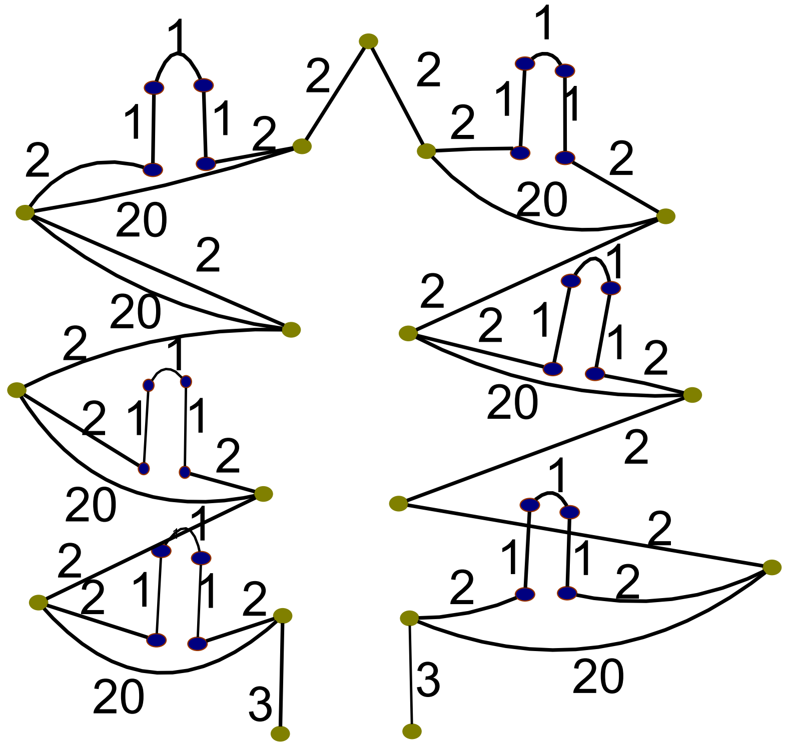

Figure 1: Poorly disguised Christmas tree.