## Homework 12

Algorithms on Directed Graphs, Winter 2018/9

Due: 1.2.2018 by 16:00

**Definition.** Let G = (V, E) be a digraph. We say that G is **Eulerian** if G is (weakly) connected, and for all  $v \in V$  we have  $|\delta^+(v)| = |\delta^-(v)|$ , where  $\delta^+(v)$  and  $\delta^-(v)$  denote the sets of out- and in-edges of v, respectively. An **Eulerian circuit** is an edge-simple cycle C that contains every edge  $e \in E$  (exactly once).

**Exercise 1.** Prove that if G is Eulerian, then it contains an Eulerian circuit. **Exercise 2.** Suppose G = (V, E) is Eulerian, and let  $A, B \subseteq V$  be disjoint. Let  $\Gamma = (V, E')$  be the undirected graph formed by replacing each directed edge  $(u, v) \in E$  with the undirected edge  $\{u, v\}$  in E'.

(a) Show that if (A, B) is a partition of V (i.e,  $V = A \cup B$ ), then we have

 $|\{(u,v) \in E \mid u \in A, v \in B\}| = |\{(u,v) \in E \mid u \in B, v \in A\}|.$ 

That is, the number of edges from A to B is equal to the number of edges from B to A.

(b) Suppose that for all  $v \in V$ , the in-degree of v satisfies  $|\delta^-(v)| \leq \Delta$ . Suppose that in  $\Gamma$ , there are  $(\Delta + 1)k + 1$  vertex disjoint (undirected) paths from A to B. Prove that in G, there are k + 1 directed vertex disjoint paths from A to B. (Do not assume here that A and B form a partition of V—they can be arbitrary disjoint sets.)

**Definition.** Let G = (V, E) be a digraph, and let  $A, B \subseteq V$ .

- We say (A, B) is a *separation* of order  $k = |A \cap B|$  if there is no (directed) edge  $(u, v) \in E$  with  $u \in A \setminus B$  and  $v \in B \setminus A$ .
- A subset X ⊂ V is node well-linked (NWL) if for all A, B ⊆ X with |A| = |B| there are k = |A| = |B| vertex disjoint paths from A to B in G.
- A subset  $X \subset V$  is *edge well-linked* (EWL) if for all  $A, B \subseteq X$  with |A| = |B| there are k = |A| = |B| edge disjoint paths from A to B in G.

- For  $\alpha \in [0, 1]$ , a subset  $X \subseteq V$  is  $\alpha$ -NWL if every separation (A, B) has order at least  $\alpha \min\{|X \cap A|, |X \cap B|\}$ .
- For  $\alpha \in [0, 1]$ , a subset  $X \subseteq V$  is  $\alpha$ -EWL if for every partition  $V = A \cup B$  we have  $|\delta^+(A)| \ge \alpha \min\{|X \cap A|, |X \cap B|\}.$

**Exercise 3.** Let G = (V, E) be a digraph and  $X \subseteq V$ .

- 1. Show that if X is 1-NWL then X is NWL.
- 2. Show that if X is 1-EWL then X is EWL.

**Definition.** Let G = (V, E) be a digraph and  $S \subset V$ . The *edge expansion* of S is defined to be

$$\alpha_S(G) = \frac{|\delta^+(S)|}{\min\{|S|, |V \setminus S|\}}.$$

The *edge expansion* of G is

$$\alpha(G) = \min_{S \subset V} \alpha_S(G).$$

**Exercise 4.** Let G = (V, E) be a digraph.

- (a) Prove that if G has edge expansion  $\alpha = \alpha(G)$  then V is  $\alpha$ -EWL.
- (b) Prove that if G is  $\alpha$ -EWL and has maximum in-degree  $\Delta$ , then G is  $(\alpha/\Delta)$ -NWL.