max planck institut
informatik

Universität
des
Saarlandes
FR Informatik

Christoph Weidenbach **October 23, 2018**

Tutorials for "Automated Reasoning WS18/19"
Exercise sheet 1

Exercise 1.1:
Determine which of the following formulas are valid/satisfiable/unsatisfiable using propositional semantics, i.e., the definition of $\models$:

1. $(P \wedge Q) \rightarrow (P \vee Q)$
2. $(P \vee Q) \rightarrow (P \wedge Q)$
3. $\neg(P \rightarrow \neg P)$
4. $(P \vee \neg Q) \wedge \neg(\neg P \rightarrow \neg Q)$
5. $\neg(P \vee Q) \leftrightarrow (\neg P \wedge \neg Q)$

Exercise 1.2:
Prove the validity of the following formulas using $\Rightarrow_{\mathrm{T}}$.

1. $(P \rightarrow (Q \rightarrow R)) \rightarrow ((P \rightarrow Q) \rightarrow (P \rightarrow R))$
2. $(P \rightarrow Q) \rightarrow ((R \vee P) \rightarrow (R \vee Q))$

Exercise* 1.3:
Consider a satisfiable formula ϕ with $\mathcal{A} \models \phi$.

1. Prove $\Rightarrow_{\mathrm{T}}$ to be strongly complete with respect to models: if $\{(\phi)\} \Rightarrow_{\mathrm{T}}^{*} N$ and N is a normal form then there is a sequence $(\phi, \phi_1, \ldots, \phi_n) \in N$ such that $\mathcal{A} \models \phi \wedge \phi_1 \wedge \ldots \wedge \phi_n$.
2. Is $\mathcal{A}$ the only model of $\phi \wedge \phi_1 \wedge \ldots \wedge \phi_n$?

Is is not encouraged to prepare joint solutions, because we do not support joint exams.