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**Tutorials for “Decision Procedures SS19”**  
**Exercise sheet 8**

**Exercise 8.1:**

Let  $\Lambda_1 \parallel C_1 \vee L_1$  and  $\Lambda_2 \parallel C_2 \vee L_2$  be two BS clauses with simple bounds. Let  $\sigma$  be the unifier of  $L_1$  and  $\text{comp}(L_2)$ . Prove that  $(\Lambda_1, \Lambda_2 \parallel C_1 \vee C_2)\sigma$  is a BS clause with simple bounds. This might need some simplification.

**Exercise 8.2:**

For the pure BS fragment the following rules are sound, complete, and terminate:

**Superposition-BS**  $M \uplus \{N \uplus \{P(t_1, \dots, t_n), C \vee \neg P(s_1, \dots, s_n)\}\} \Rightarrow_{\text{SUPBS}} M \cup \{N \cup \{P(t_1, \dots, t_n), C \vee \neg P(s_1, \dots, s_n)\} \cup \{C\sigma\}\}$

where (i)  $\neg P(s_1, \dots, s_n)$  is selected in  $(C \vee \neg P(s_1, \dots, s_n))\sigma$  (ii)  $\sigma$  is the mgu of  $P(t_1, \dots, t_n)$  and  $P(s_1, \dots, s_n)$  (iii)  $C \vee \neg P(s_1, \dots, s_n)$  is a Horn clause

**Instantiation**  $M \uplus \{N \uplus \{C \vee A_1 \vee A_2\}\} \Rightarrow_{\text{SUPBS}} M \cup \{N \cup \{(C \vee A_1 \vee A_2)\sigma_i \mid \sigma_i = \{x \mapsto a_i\}, 1 \leq i \leq k\}\}$

where  $x$  occurs in a variable chain between  $A_1$  and  $A_2$

**Split**  $M \uplus \{N \uplus \{C_1 \vee A_1 \vee C_2 \vee A_2\}\} \Rightarrow_{\text{SUPBS}} M \cup \{N \cup \{C_1 \vee A_1\}, N \cup \{C_2 \vee A_2\}\}$

where  $\text{vars}(C_1 \vee A_1) \cap \text{vars}(C_2 \vee A_2) = \emptyset$

Discuss to what extent the rules can be turned into a sound, complete and terminating calculus for BS with simple bounds over LRA.