2.9.10 Lemma (CDCL Soundness)

In a reasonable CDCL derivation, CDCL can only terminate in two different types of final states: $(M; N; U; k; \top)$ where $M \models N$ and $(M; N; U; k; \bot)$ where N is unsatisfiable.



2.9.11 Proposition (CDCL Soundness)

The rules of the CDCL algorithm are sound: (i) if CDCL terminates from $(\epsilon; N; \emptyset; 0; \top)$ in the state $(M; N; U; k; \top)$, then N is satisfiable, (ii) if CDCL terminates from $(\epsilon; N; \emptyset; 0; \top)$ in the state $(M; N; U; k; \bot)$, then N is unsatisfiable.



2.9.12 Proposition (CDCL Strong Completeness)

The CDCL rule set is complete: for any valuation M with $M \models N$ there is a reasonable sequence of rule applications generating $(M'; N; U; k; \top)$ as a final state, where M and M' only differ in the order of literals.



2.9.13 Proposition (CDCL Termination)

Assume the algorithm CDCL with all rules except Restart and Forget is applied using a reasonable strategy. Then it terminates in a state (M; N; U; k; D) with $D \in \{\top, \bot\}$.



Application		
${\sf System} + {\sf Problem}$		
System		
Algorithm + Implementation		
Algorithm		
Calculus + Strategy		
Calculus		
Logic + States + Rules		
Logic		
Syntax + Semantics		



```
1 Algorithm: 5 CDCL(S)
   Input : An initial state (\epsilon; N; \emptyset; 0; \top).
   Output: A final state S = (M; N; U; k; \top) or S = (M; N; U; k; \bot)
 2 while (any rule applicable) do
       ifrule (Conflict(S)) then
 3
           while (Skip(S) \parallel Resolve(S)) do
 4
              update VSIDS on resolved literals;
 5
           update VSIDS on learned clause, Backtrack(S);
 6
           if (forget heuristic) then
              Forget(S), Restart(S);
 8
          else
              if (restart heuristic) then
10
                  Restart(S):
11
       else
12
           ifrule (! Propagate(S)) then
13
              Decide(S) literal with max. VSIDS score: _{November 5, 2020}
```

```
Propagate (M; N; U; k; \top) \Rightarrow_{CDCL} (ML^{C \lor L}; N; U; k; \top) provided C \lor L \in (N \cup U), M \models \neg C, and L is undefined in M
```

Conflict
$$(M; N; U; k; \top) \Rightarrow_{\mathsf{CDCL}} (M; N; U; k; D)$$
 provided $D \in (N \cup U)$ and $M \models \neg D$



- data structures: clauses, trail, and the rules
- heuristics: decision literal, forget, restart
- space efficiency: forget
- quality: restarts
- special cases





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Data Structures

Idea: Select two literals from each clause for indexing.



Data Structures

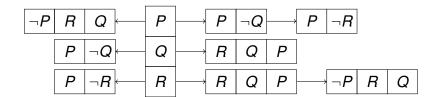
Idea: Select two literals from each clause for indexing.

2.10.1 Invariant (2-Watched Literal Indexing)

If one of the watched literals is false and the other watched literal is not true, then all other literals of the clause are false.



$$N = \{P \lor \neg R, P \lor \neg Q, R \lor Q \lor P, \neg P \lor R \lor Q\}$$





VSIDS: Variable State Independent Decaying Sum

- each propositional variable has a positive score, initially 0
- decide the variable with maximal score, remember sign (phase saving)
- increment the score of variables involved in resolution by b
- increment the score of variables in learned clauses by b
- initially b > 0
- at Backtrack set b := c * b where 2 >> c > 1, i.e., $b_n = c^n * b$
- take care of overflows, i.e., rescale from time to time
- sometimes pick a variable randomly





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- fix a limit d on the number of learned clauses
- if more than |U| > d start forgetting
- remove redundant clauses
- sort the learned clauses according to a score
- typical elements of the score are clause length, the VSIDS score, dependency on decisions
- remove the k% clauses with minimal score from U
- d := d + e for some e, e >> 1
- do a Restart





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Restart

- after forgetting do a restart
- if a unit is learned do a restart
- restart often at the beginning of a run
- classics: Luby sequence 1, 1, 2, 1, 1, 2, 4, ... $(u_1, v_1) := (1, 1),$ $(u_{n+1}, v_{n+1}) := ((u_n \& u_n) = v_n?(u_n + 1, 1) : (u_n, 2 * v_n))$

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Memory Matters: SPASS-SATT

Forget-Start	800	108800
Restarts	412	369
Conflicts	153640	133403
Decisions	184034	159005
Propagations	17770298	15544812
Time	11	23
Memory	16	41



Propositional Logic Calculi

- 1. Tableau: classics, natural from the semantics
- 2. Resolution: classics, first-order, prepares for CDCL
- 3. CDCL: current prime calculus for propositional logic
- 4. Superposition: first-order, prepares for first-order

